#### The Smallest Hard Trees

#### Florian Starke

joint work with Manuel Bodirsky, Jakub Bulín, and Michael Wernthaler







The CSP of G is the following Problem:

Input: H

Output: 3 homomorphism H-> G?

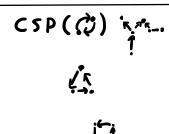
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Output: 3 homomorphism H-> G?



The CSP of G is the following Problem:

Input: He finite digraphs

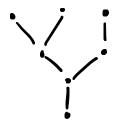
Output: 3 homomorphism H-> 6?

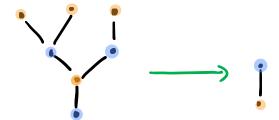
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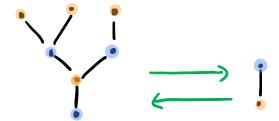
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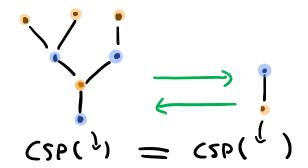
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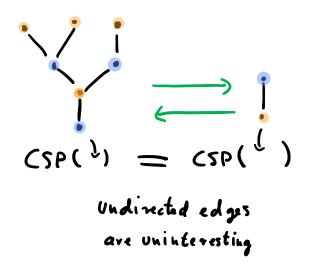
H-) is iff H contains

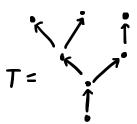


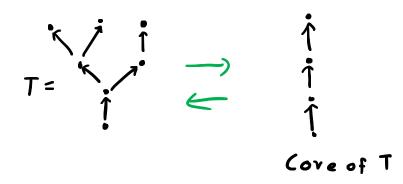


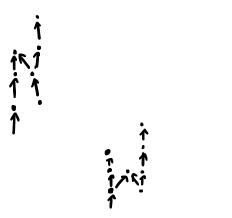














# History of NP-Hard Trees

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# Goal: assuming P+NP

Find the smallest NP-hard Tree.

TFAE (assuming P+NP)

- · CSP(T) is NP-hard
- T can pp-construct of f(area) = f(rare)
- · Thas ho Siggers polymorphism

# TFAE (ASSUMING P+NP)

• T can pp-construct of area)
• T has no Siggers-polymorphism

(can be computed by testing

Taveau rave

contract tuples that should have the same image

# TFAE (assuming P+NP)

· CSP(T) is NP-hard • T can pp-construct = f(area) · T has no Siggers - polymorphism can be computed by testing 204 THE TAVERS THAT Should

# TFAE (assuming P+NP)

- · CSP(T) is NP-hard
- T can pp-construct = f(area) = f(rare)
- · Thas ho Siggers polymorphism
- T has no Kearnes-Marković-Mckmzie-Polymorphisus

  P(abb)= q(baa)=q(aab)

  p(aba)=q(aba)

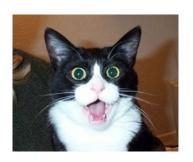
#### TFAE (assuming PANP)

- · CSP(T) is NP-hard
- T can pp-construct of f(area) = f(rare)
- · Thas ho Siggers polymorphism
- P(abb)= q(baa)=q(aab) 2.203 = 16000
  P(aba)=q(aba) huch before

but still not easy

n	trees
1	1
2	1
3	3
4	8
5	27
6	91
7	350
8	1376
9	5743
10	24635

n	trees
1	1
2	1
3	3
4	8
5	27
6	91
7	350
8	1376
9	5743
10	24635
11	108968
12	492180
13	2266502
14	10598452
15	50235931
16	240872654
17	1166732814
18	5702001435
19	28088787314
20	139354922608



testing 100 Evers per second this would take 317 years

n	trees	cores
1	1	1
2	1	1
3	3	1
4	8	1
5	27	1
6	91	2
7	350	3
8	1376	7
9	5743	15
10	24635	36
11	108968	85
12	492180	226
13	2266502	578
14	10598452	1569
15	50235931	4243
16	240872654	11848
17	1166732814	33104
18	5702001435	94221
19	28088787314	269455
20	139354922608	779268

testing 100 trees per second this would take 2 hours

much better but still not easy

**Theorem** Let  $\mathbb{T}$  be a finite tree. Then the following are equivalent.

- 1.  $\mathbb{T}$  is a core;
- 2.  $\operatorname{End}(\mathbb{T}) = \{ \operatorname{id}_T \};$
- 3.  $AC_{\mathbb{T}}(\mathbb{T})$  terminates such that the list for each vertex contains a single element.

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for each vertex contains a single element.

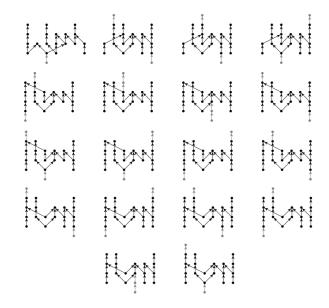


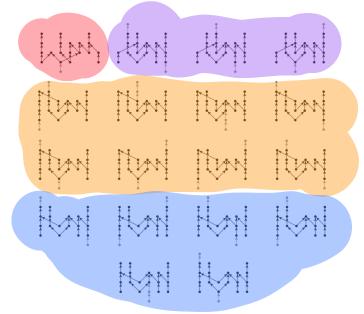
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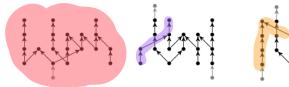


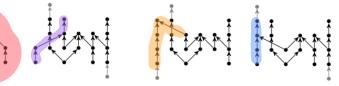
foreach core tree T

IF T has no kMM-polymorphisms than return T

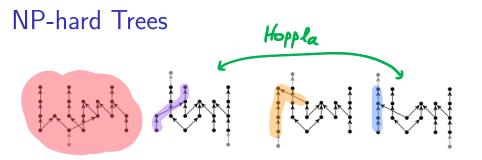




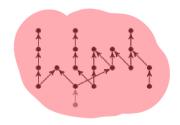


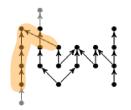


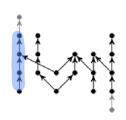
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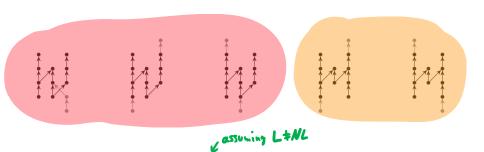




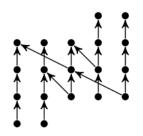


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			Taccoming DL MD





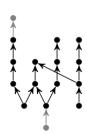
smallest NL-hard trees
(12 vertices)



smallest tree not solved by AC
(19 vertices)

# Open Problem

# Open Problem



- in P
- NL-hard
- Smallest tree without majority polymorphism
- has kearnes-kiss-chain of length 5 (no kk-clain=) P-hand)
- has no Jóhnsson-chain of Lugth 1000 (J-chain =) in NL)

(16 vertices)

# Paper: The Smallest Hard Trees

https://arxiv.org/abs/2205.07528