

Adjunct to the polymorphism functor

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Endopolymorphisms

Given a relational structure \mathbf{A} .

We say that $f: A^n \rightarrow A$ is a **polymorphism** of \mathbf{A} if one of the following equivalent conditions is satisfied:

- ▶ f is a homomorphism from \mathbf{A}^n to \mathbf{A} ,
- ▶ for each relation $R^{\mathbf{A}}$ and all tuples $\mathbf{a}_1, \dots, \mathbf{a}_n \in R^{\mathbf{A}}$ we have

$$f(\mathbf{a}_1, \dots, \mathbf{a}_n) \in R^{\mathbf{A}},$$

- ▶ each relation $R^{\mathbf{A}}$ is a subuniverse of $(A; f)^k$ where k is the arity of R .

The set **Pol**(\mathbf{A}) of all (endo)polymorphisms of \mathbf{A} is a clone.

Polymorphisms

Given relational structures \mathbf{A} and \mathbf{B} that share a signature.

We say that $f: A^n \rightarrow B$ is a **polymorphism** from \mathbf{A} to \mathbf{B} if one of the following equivalent conditions is satisfied:

- ▶ f is a homomorphism from \mathbf{A}^n to \mathbf{B} ,
- ▶ for each relation $R^{\mathbf{A}}$ and all tuples $\mathbf{a}_1, \dots, \mathbf{a}_n \in R^{\mathbf{A}}$ we have

$$f(\mathbf{a}_1, \dots, \mathbf{a}_n) \in R^{\mathbf{B}},$$

- ▶ each relational pair $(R^{\mathbf{A}}, R^{\mathbf{B}})$ is a subuniverse of $(A, B; f)^k$ where k is the arity of R .

The set $\text{Pol}(\mathbf{A}, \mathbf{B})$ of all polymorphisms from \mathbf{A} to \mathbf{B} is **not** a clone, but it is closed under taking minors.

Minor closed sets

a.k.a. clonoids

Let $f: A^n \rightarrow B$ be a function. Any function g of the form

$$g(x_1, \dots, x_m) = f(x_{e(1)}, \dots, x_{e(n)}).$$

for some $e: [n] \rightarrow [m]$ is called a **minor** of f .

Theorem (Pippenger, 2002; Brakiensiek, Guruswami, 2016)

For all finite sets A, B and every minor closed set $\mathcal{A} \subseteq \mathcal{O}(A, B)$ there exist relational structures \mathbf{A} and \mathbf{B} such that $\text{Pol}(\mathbf{A}, \mathbf{B}) = \mathcal{A}$.

$(\mathcal{O}(A, B) = \{f \mid f: A^n \rightarrow B, n \in \mathbb{N}\})$

Minor preserving maps

a.k.a. h1 homomorphisms, clonoid homomorphisms

Let \mathcal{A} and \mathcal{B} be minor closed. A map $\xi: \mathcal{A} \rightarrow \mathcal{B}$ that preserves arities is **minor preserving** if for each $f \in \mathcal{A}^{(n)}$ and each $e: [n] \rightarrow [m]$ we have

$$\xi(f(\pi_{e(1)}^m, \dots, \pi_{e(n)}^m)) \approx \xi(f)(\pi_{e(1)}^m, \dots, \pi_{e(n)}^m).$$

Theorem (Barto, O, Pinsker, 2017)

If \mathbf{A} and \mathbf{B} are finite (ω -categorical) structures such that there is a (uniformly continuous) minor preserving map $\xi: \text{Pol}(\mathbf{A}) \rightarrow \text{Pol}(\mathbf{B})$, then $\text{CSP}(\mathbf{A}) \geq_L \text{CSP}(\mathbf{B})$.

The polymorphism functor

Note. Even if \mathbf{A} and \mathbf{B} are homomorphically equivalent, there is no clone homomorphism from $\text{Pol}(\mathbf{A})$ to $\text{Pol}(\mathbf{B})$.

Fix a relational signature, and let $\mathbf{A}, \mathbf{B}, \mathbf{A}', \mathbf{B}'$ be structures in this signatures, and $a: \mathbf{A}' \rightarrow \mathbf{A}, b: \mathbf{B} \rightarrow \mathbf{B}'$ homomorphisms.

$$\begin{array}{ccc} \mathbf{A}^n & \xleftarrow{a} & (\mathbf{A}')^n \\ \downarrow f & & \downarrow \xi_{a,b}(f) \\ \mathbf{B} & \xrightarrow{b} & \mathbf{B}' \end{array}$$

$$\xi_{a,b}(f)(x_1, \dots, x_n) = bf(a(x_1), \dots, a(x_n))$$

Note. $\text{Pol}(\mathbf{A}', \mathbf{B}')$ contains a **reflection** of $\text{Pol}(\mathbf{A}, \mathbf{B})$.

The adjunct

For each structure \mathbf{A} , $\text{Pol}(\mathbf{A}, -)$ has a left adjunct $\text{Free}(-, \mathbf{A})$.

Given a minor closed set \mathcal{C} and a relational structure \mathbf{A} . We define the **free structure** ('free action of \mathcal{C} on \mathbf{A} ') \mathbf{F} :

- ▶ Let $F = \{f(a_1, \dots, a_n) : n \in \mathbb{N}, f \in \mathcal{C}^{(n)}, a_1, \dots, a_n \in A\} / \approx$
- ▶ for a relation R , we define
$$R^{\mathbf{F}} = \{f(\mathbf{a}_1, \dots, \mathbf{a}_n) : n \in \mathbb{N}, f \in \mathcal{C}^{(n)}, \mathbf{a}_1, \dots, \mathbf{a}_n \in R^{\mathbf{A}}\}.$$

We use the symbol $\text{Free}(\mathcal{C}, \mathbf{A})$ for \mathbf{F} .

Observation

For each \mathbf{A}, \mathbf{B} , relational structures and \mathcal{C} , minor closed, there is a natural isomorphism $\eta_{\mathcal{C}, \mathbf{A}, \mathbf{B}}$:

$$\{h: \text{Free}(\mathcal{C}, \mathbf{A}) \rightarrow \mathbf{B}\} \stackrel{\eta_{\mathcal{C}, \mathbf{A}, \mathbf{B}}}{\simeq} \{\xi: \mathcal{C} \rightarrow \text{Pol}(\mathbf{A}, \mathbf{B})\}.$$

Linear Mal'cev conditions

A **coloring** of a minor closed set \mathcal{C} by \mathbf{A} is a homomorphism $c: \text{Free}(\mathcal{C}, \mathbf{A}) \rightarrow \mathbf{A}$.

A coloring of a clone is **strong** if $c(a) = a$ for all $a \in A$.

Note. \mathcal{C} is strongly colorable by \mathbf{A} iff there is a minor preserving map from \mathcal{C} to $\text{Pol}(\mathbf{A})$ that maps the identity map to itself.

Theorem (Sequeira, Greenwell & Lovász, ...)

Let \mathcal{C} be a clone.

- ▶ \mathcal{C} has a Mal'cev term iff it is not strongly colorable by $\mathbf{L} = (\{0, 1, 2\}, 01|2, 0|12)$.
- ▶ \mathcal{C} satisfies a non-trivial Mal'cev condition iff it is not colorable by \mathbf{K}_3 .

Deciding triviality of Mal'cev conditions

Label Cover

Mal'cev condition is **linear** if it contains only identities of the form

$$f(x_{\pi(1)}, \dots, x_{\pi(m)}) \approx g(x_{\sigma(1)}, \dots, x_{\sigma(n)}), \quad \text{or} \quad f(x_{\pi(1)}, \dots, x_{\pi(m)}) \approx x_1.$$

Corollary

Given a linear Mal'cev condition Σ of arity at most N .

- ▶ *For each N , deciding whether Σ implies the Mal'cev term is solvable in Ptime.*
- ▶ *For each $N \geq 6$, deciding whether Σ is trivial is NP-complete.*

Proof. Construct $\text{Free}(\Sigma, \mathbf{A})$ in Ptime, then decide existence of a homomorphism $\text{Free}(\Sigma, \mathbf{A}) \rightarrow \mathbf{A}$ by $\text{CSP}(\mathbf{A})$. ■

Promise constraint satisfaction

Fix two finite relational structures \mathbf{A} and \mathbf{B} with the same finite signature. $\text{PCSP}(\mathbf{A}, \mathbf{B})$ is the following problem: Given a structure \mathbf{Q} in the common language, output

- ▶ YES if \mathbf{Q} maps homomorphically into \mathbf{A} ,
- ▶ NO if \mathbf{Q} does not map homomorphically into \mathbf{B} .

Note. $\text{CSP}(\mathbf{A}) \equiv \text{PCSP}(\mathbf{A}, \mathbf{A})$ and $\text{PCSP}(\mathbf{A}, \mathbf{B}) \leq \text{CSP}(\mathbf{A}), \text{CSP}(\mathbf{B})$.

Gap Label Cover

Fix \mathbf{A}, \mathbf{B} . Given a minor closed set \mathcal{C} , we know

- ▶ $\mathcal{C} \rightarrow \text{Pol}(\mathbf{A}, \mathbf{A})$ iff $\text{Free}(\mathcal{C}, \mathbf{A}) \rightarrow \mathbf{A}$, and
- ▶ $\mathcal{C} \not\rightarrow \text{Pol}(\mathbf{A}, \mathbf{B})$ iff $\text{Free}(\mathcal{C}, \mathbf{A}) \not\rightarrow \mathbf{B}$.

Gap Mal'cev Sat. Fix two minor closed sets \mathcal{A} and \mathcal{B} . Given Σ of maximal arity N , output

- ▶ YES if Σ is satisfied in \mathcal{A} ,
- ▶ NO if Σ is not satisfied in \mathcal{B} .

Denote this problem $\text{GMS}_{\mathcal{A}, \mathcal{B}}(N)$.

Reduction between LC and PCSP

Theorem (O, Bulin, 2017*)

Let \mathbf{A} and \mathbf{B} be relational structures, $\mathcal{A} = \text{Pol}(\mathbf{A}, \mathbf{A})$, and $\text{Pol}(\mathbf{A}, \mathbf{B})$.
Then

- ▶ For all $N > 0$, $\text{GMS}_{\mathcal{A}, \mathcal{B}}(N) \leq_L \text{PCSP}(\mathbf{A}, \mathbf{B})$.
- ▶ There exists $N > 0$ s.t. $\text{PCSP}(\mathbf{A}, \mathbf{B}) \leq_L \text{GMS}_{\mathcal{A}, \mathcal{B}}(N)$.

Corollary

The complexity of $\text{PCSP}(\mathbf{A}, \mathbf{B})$ depends only on minor Mal'cev conditions satisfied by $\text{Pol}(\mathbf{A}, \mathbf{B})$.

Thanks! ■